



Reconfiguration graphs induced by rainbow spanning trees

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Joint work with

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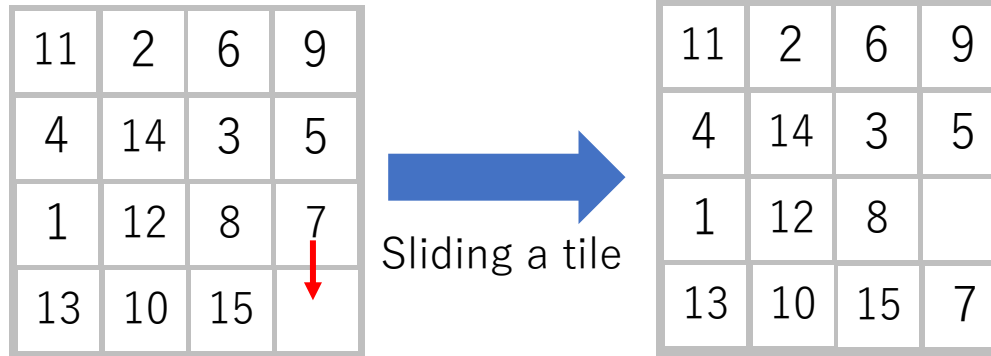
Combinatorial reconfiguration

Combinatorial reconfiguration is an algorithmic concept that provide mathematical models and analysis for “transformations over state spaces”.

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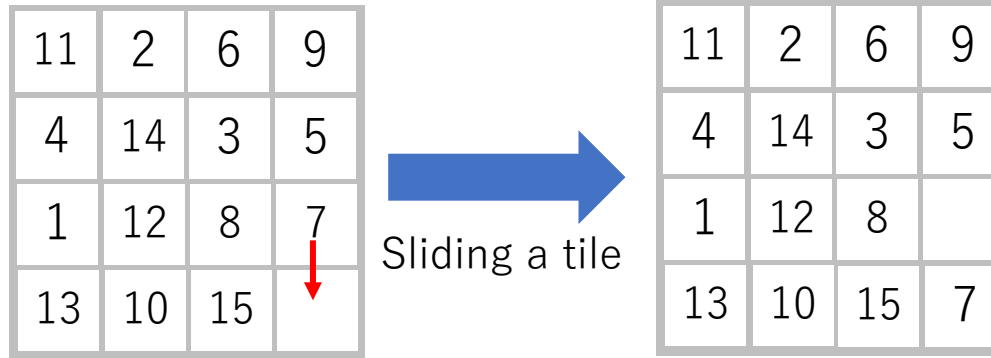
Object : 15-puzzle
Operation : Sliding tiles



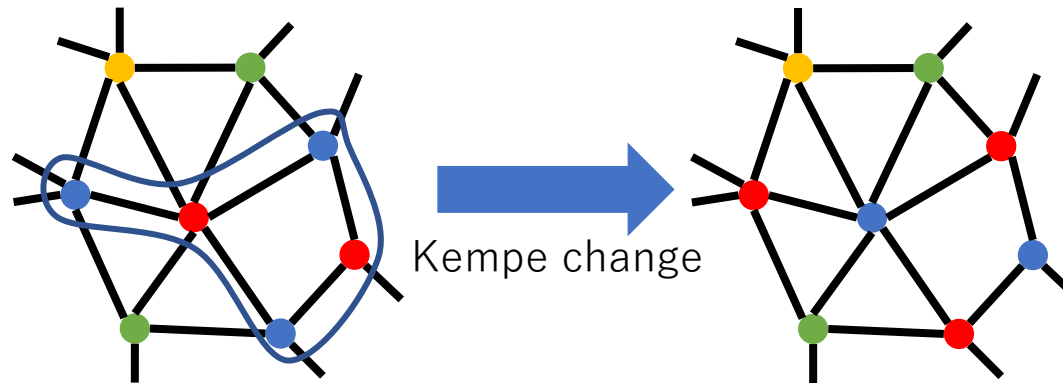
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Object : Graph Coloring
Operation : Kempe change



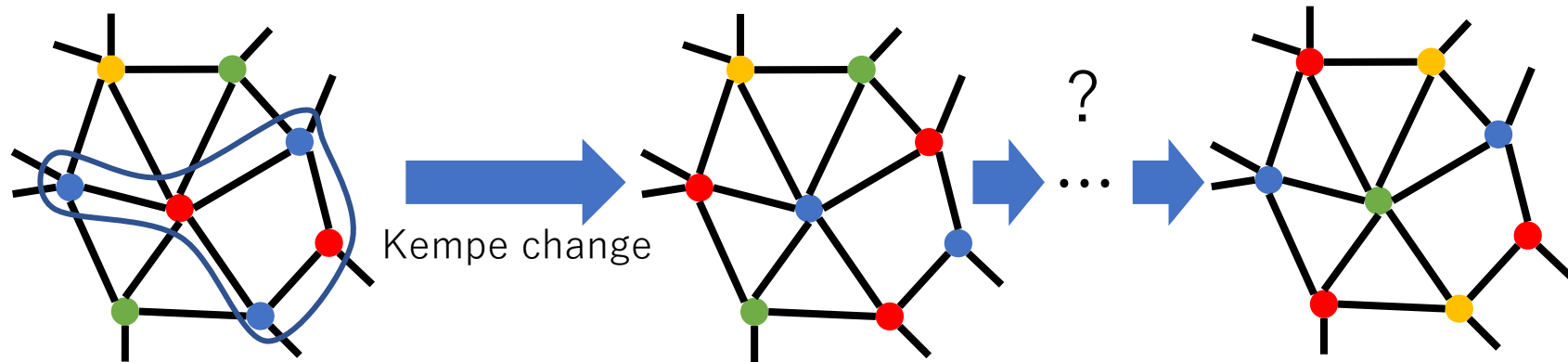
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Reachability problem: For two inputs, we are asked to determine whether or not we can transform one into the other by a prescribed operation.

Edge-colored graph and rainbow

Definition (edge-colored graph and rainbow)

- An **edge-colored graph** is a graph with an edge coloring (not necessarily proper coloring).
- An edge-colored graph is **rainbow** if no two edges have the same color.



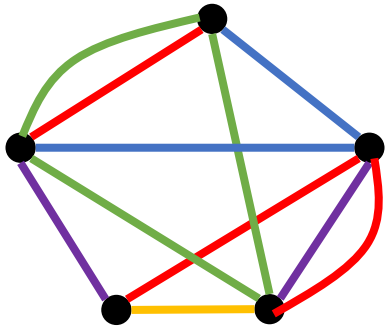
Edges in the same color
may share an end-vertex

Rainbow

Rainbow spanning tree reconfiguration

Object : Edge-colored multigraph Operation : Edge flip

Rule : All intermediate results remain rainbow spanning trees

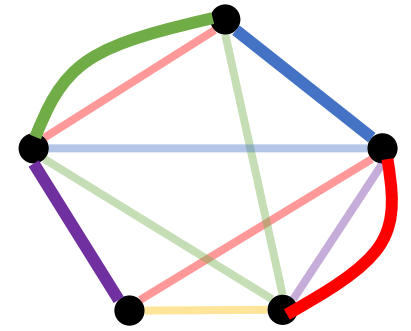
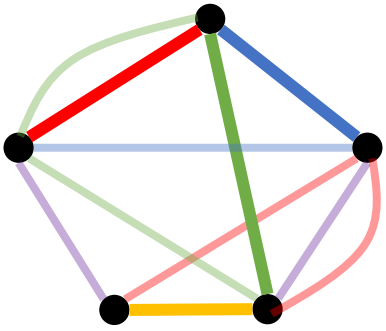


Rainbow spanning tree reconfiguration

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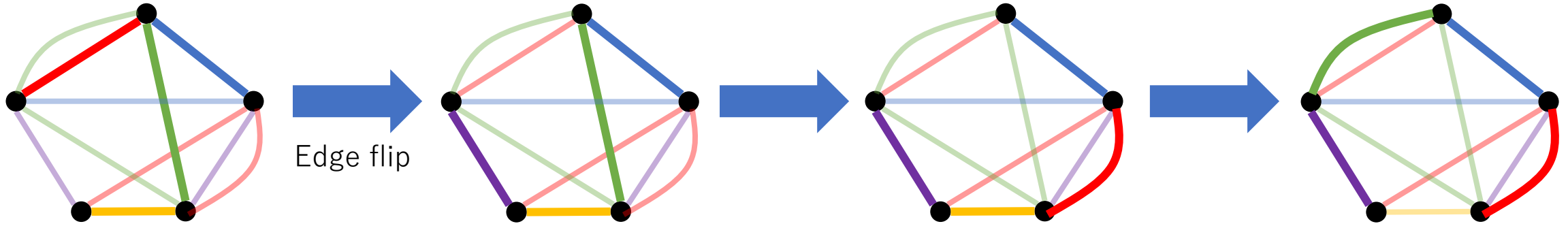


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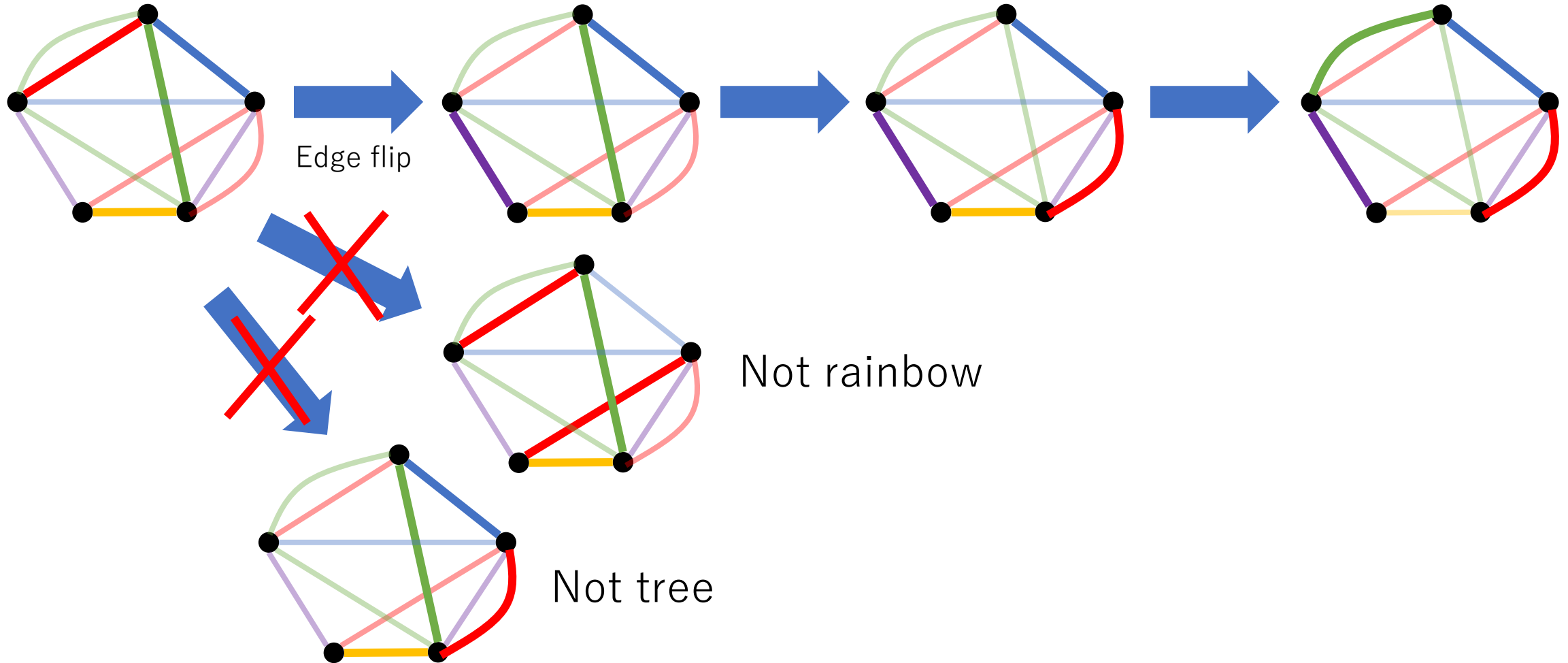


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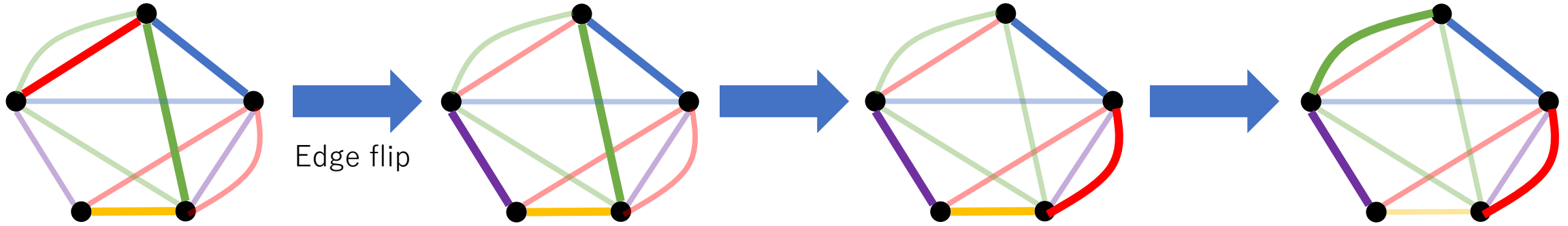


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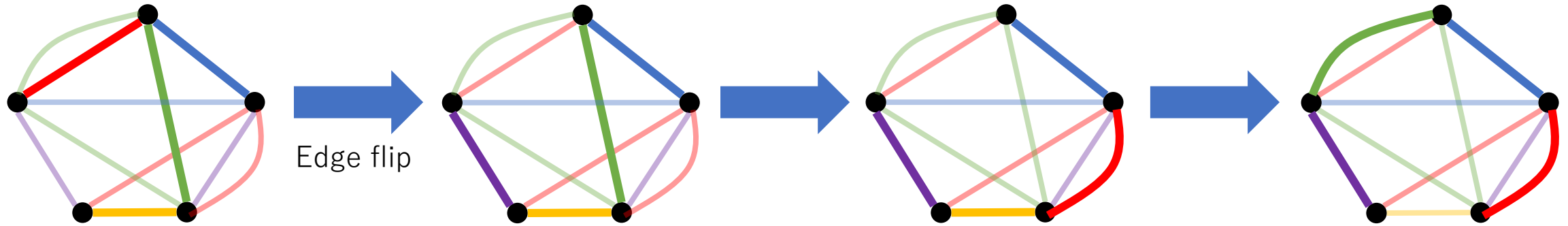
Definition (Rainbow spanning tree reconfiguration)

T, T' : Rainbow spanning trees of an edge-colored multigraph G

A **reconfiguration sequence** between T and T' is a sequence of rainbow spanning trees (T_0, T_1, \dots, T_k) in G with $T_0 = T$ and $T_k = T'$ s.t.

T_{i+1} is obtained from T_i by edge flip i.e. $T_{i+1} = T_i - e + f$.

Rainbow spanning tree reconfiguration



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Problem (Rainbow spanning tree reconfiguration problem)

Is there a polynomial-time algorithm for the following decision problem?

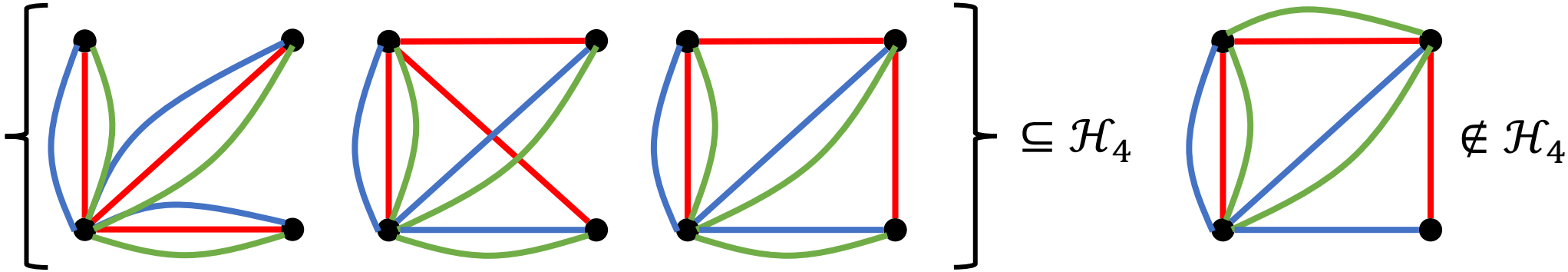
Input : edge colored multigraph G , rainbow spanning trees T, T' of G

Output : whether there is a reconfiguration sequence between T and T'

What are conditions for G to always result in “yes”?
How about **reconfiguration** graph?

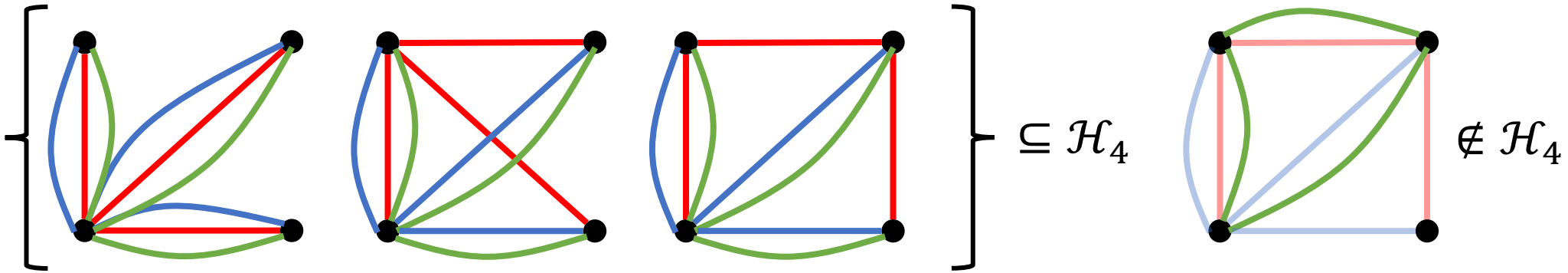
Special edge-colored graph

\mathcal{H}_n : the set of edge-colored graphs with n vertices satisfying that edges colored with a color c induce a connected spanning graph for each color c in the graph



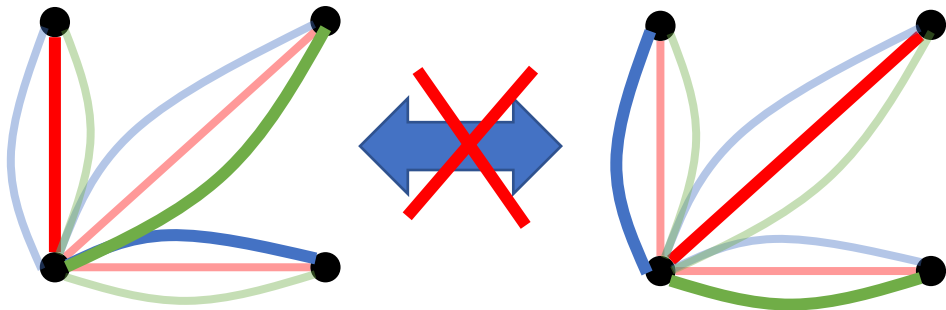
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Is there reconf. sequence between any two rainbow spanning trees in $G \in \mathcal{H}_n$?

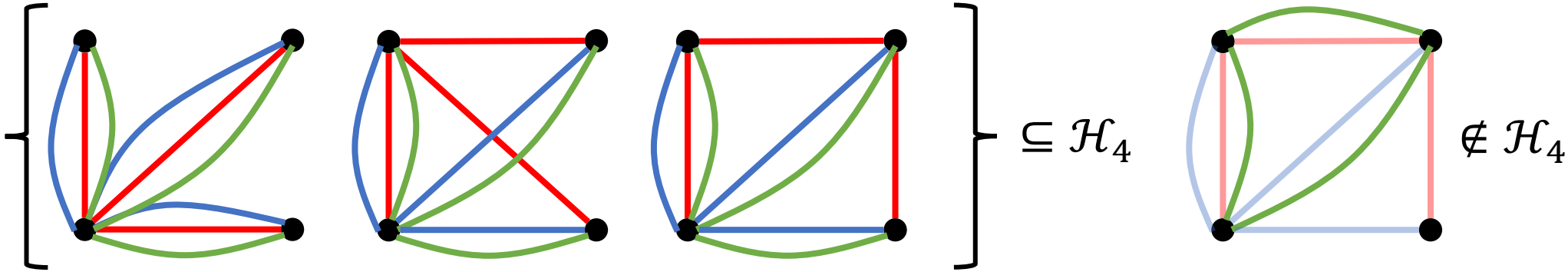
\Rightarrow No



Number of colors is too few...

Special edge-colored graph

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Theorem 1 (M and Yamaguchi)

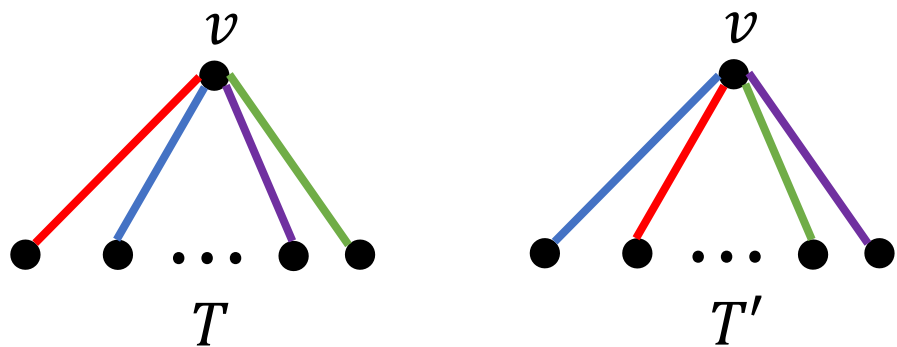
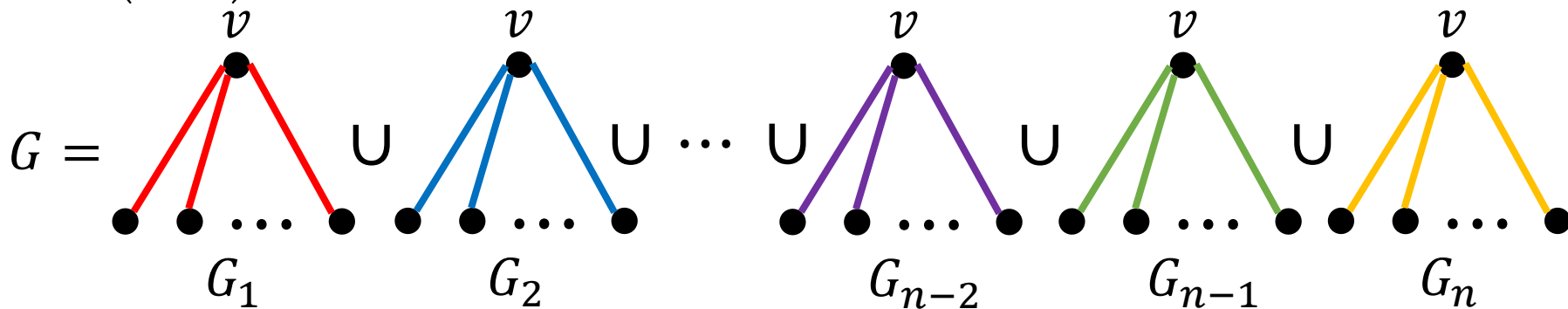
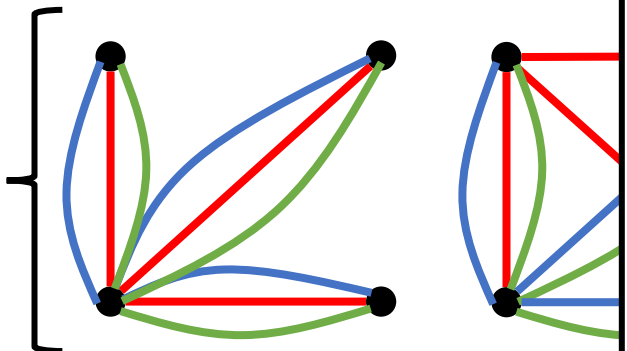
For any two rainbow spanning trees in an edge-colored graph $G \in \mathcal{H}_n$ having at least n colors, there is a reconfiguration sequence between them and the length of the sequence is at most $\frac{3}{2}(n - 1)$.

The bound is tight.

Special edge

n (odd) : number of vertices of G

\mathcal{H}_n : the set of edges with a color c induced



Theorem 1 (M and Y)
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least n colors, there
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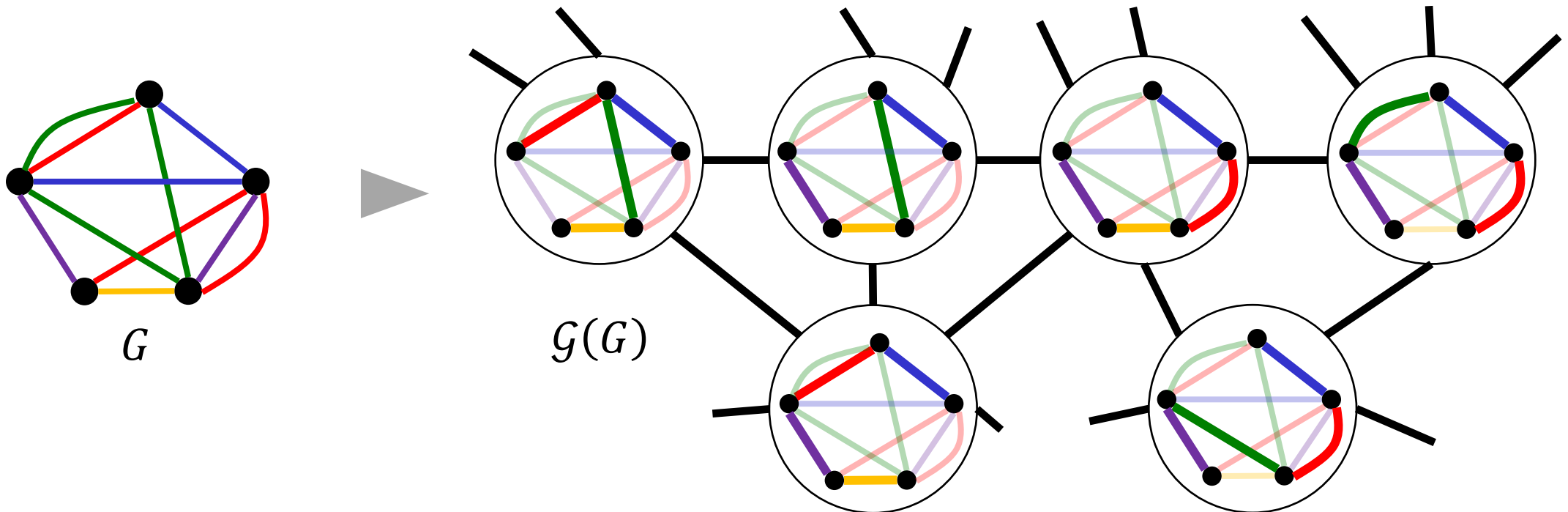
Reconfiguration graphs

Definition (Reconfiguration graph induced by rainbow spanning trees)

For an edge-colored multigraph G , **the reconfiguration graph** $\mathcal{G}(G)$ induced by rainbow spanning trees of G as follows:

$$V(\mathcal{G}(G)) = \{T : \text{rainbow spanning tree } T \text{ in } G\},$$

two rainbow spanning trees are adjacent in $\mathcal{G}(G) \Leftrightarrow$
one is obtained from the other by exchanging a single edge.



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Theorem 1 (rephrase)

For an edge-colored graph $G \in \mathcal{H}_n$ having at least n colors, the diameter of $\mathcal{G}(G)$ is at most $\frac{3}{2}(n - 1)$.

What other properties might the reconfiguration graphs have?

Reconfiguration graphs

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Definition (Reconfiguration graph induced by spanning trees)

For a multigraph G , **the reconfiguration graph** $\mathcal{G}'(G)$ induced by spanning trees of G as follows:

$$V(\mathcal{G}'(G)) = \{T : \text{spanning tree } T \text{ in } G\},$$

two spanning trees are adjacent in $\mathcal{G}'(G) \Leftrightarrow$
one is obtained from the other by exchanging a single edge.

Theorem A (Holzmann and Harary, 1972)

For a multigraph G , $\mathcal{G}'(G)$ is hamiltonian.

Reconfiguration graphs

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Problem 1

For any edge-colored multigraph $G \in \mathcal{H}_n$ having at least $n(\geq 3)$ colors, is $\mathcal{G}(G)$ hamiltonian?

Theorem 2 (M and Yamaguchi)

For any edge-colored multigraph $G \in \mathcal{H}_n$ having at least n colors, $\mathcal{G}(G)$ is 2-connected.

Problem 2

For any edge-colored multigraph $G \in \mathcal{H}_n$ having at least n colors, is $\mathcal{G}(G)$ 1-tough?

Conclusion

Thank you very much!!

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For an edge-colored graph $G \in \mathcal{H}_n$ having at least n colors, the diameter of $\mathcal{G}(G)$ is at most $\frac{3}{2}(n - 1)$.

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